

Modeling and Analyzing MAPE-K Feedback Loops for Self-adaptation

Paolo Arcaini

Dept. of Distributed and Dependable Systems, Charles University in
Prague, Czech Republic -- arcai@fd3s.mff.cuni.cz



Elvinia Riccobene

Dept. of Computer Science - University of Milan, Italy --
elvinia.riccobene@unimi.it



Patrizia Scandurra

Dept. of Management, Information and Production Engineering,
University of Bergamo, Italy -- patrizia.scandurra@unibg.it



Problem statement

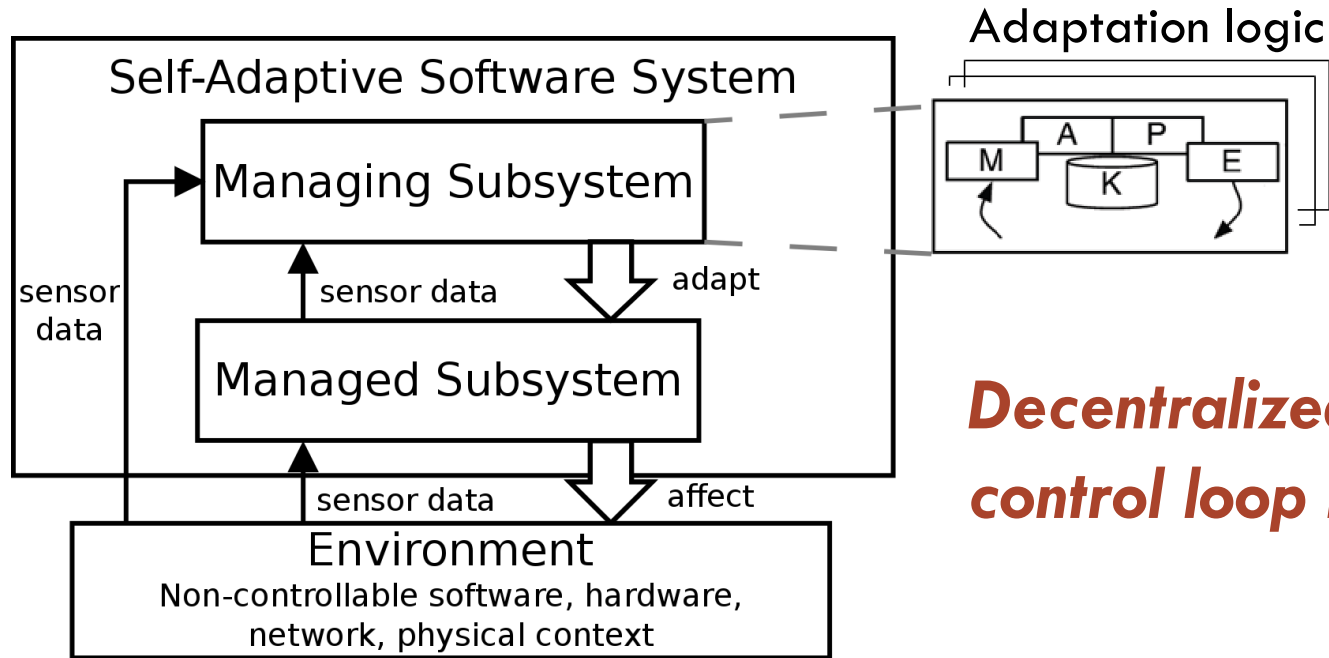
- ▶ **Self-Adaptation (SA)** is a promising approach to deal with the complexity, uncertainty and dynamicity of modern software systems
- ▶ The **MAPE-K (Monitor-Analyze-Plan-Execute over a shared Knowledge)** feedback loop is a well-known *control model* for autonomic and self-adaptive systems
- ▶ **Formal methods** for specifying and reasoning about self-adaptive systems' behavior **are highly demanded**
 - A study (reference [34] in the paper) shows the number of works that employ formal methods in self-adaptive systems are low
- ▶ Our **proposal**:
 - A **formal framework for modeling, validating, and verifying self-adaptive systems with multiple interactive MAPE-K loops**
 - based on the formal method **Abstract State Machines** and **model-checking** techniques

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Outline

- ▶ *Decentralized MAPE-K control loops*: reference model for Self-Adaptation
- ▶ Background on *Abstract State Machines (ASM)*
- ▶ *Self-adaptive ASMs*: enhanced ASM constructs and patterns to model self-adaptive behavior
- ▶ Tool-supported formal analysis techniques
- ▶ Conclusions and future work

Reference model for Self-Adaptation



Decentralized MAPE-K control loop model

- ▶ **MAPE-K (Monitor-Analyse-Plan-Execute components over a shared Knowledge)** : well known architectural solution to realize the control loop of a self-adaptive system
 - J. O. Kephart and D. M. Chess. The vision of autonomic computing. IEEE Computer, 36(1):41-50, 2003
- ▶ **Separation of concerns:** a set of interacting MAPE loops, one per each adaptation concern
- ▶ **Decentralization:** MAPE computations may be decentralized throughout multiple MAPE loops
 - They need to be coordinated to avoid conflicts!

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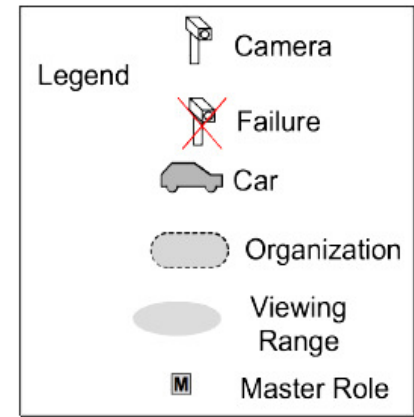
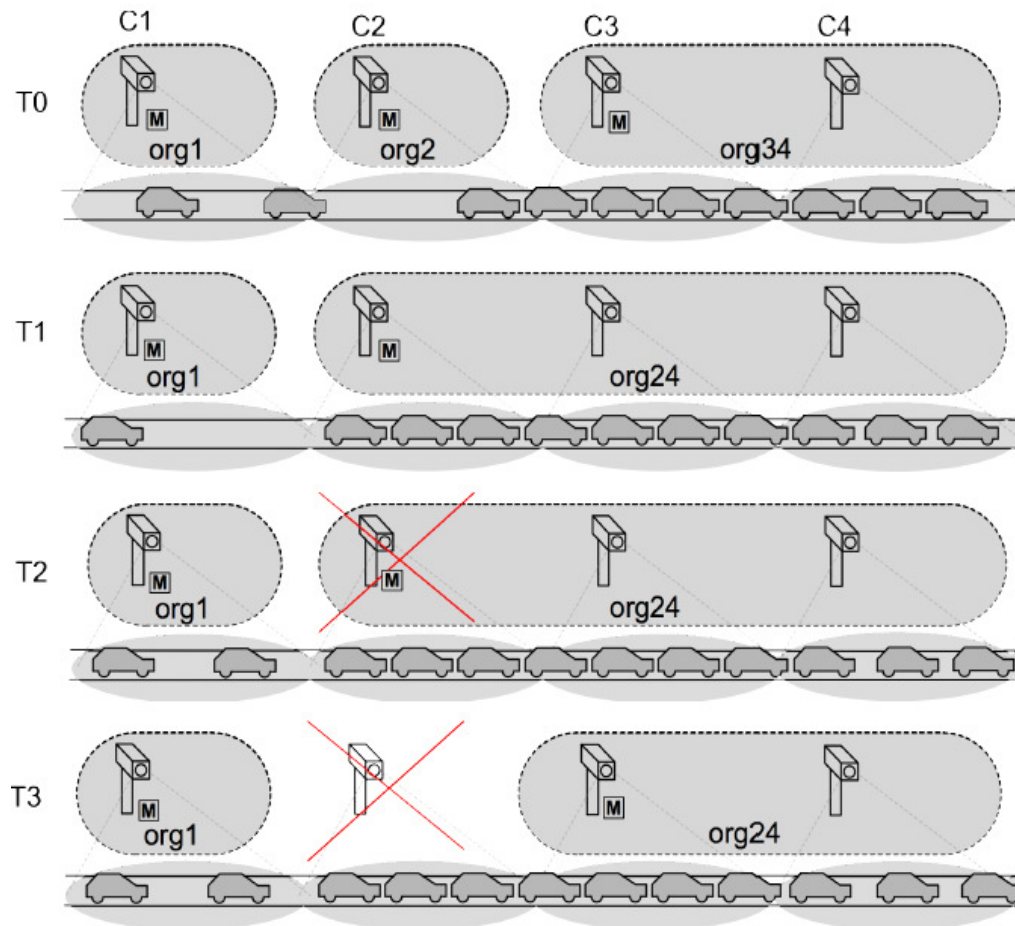
Running case study: Traffic Monitoring application

(inspired by *)

Intelligent cameras collaborate in *master/slaves organizations* to monitor and aggregate useful data whenever the traffic jam enters/leaves their viewing range

FLEXIBILITY

adaptation
concern



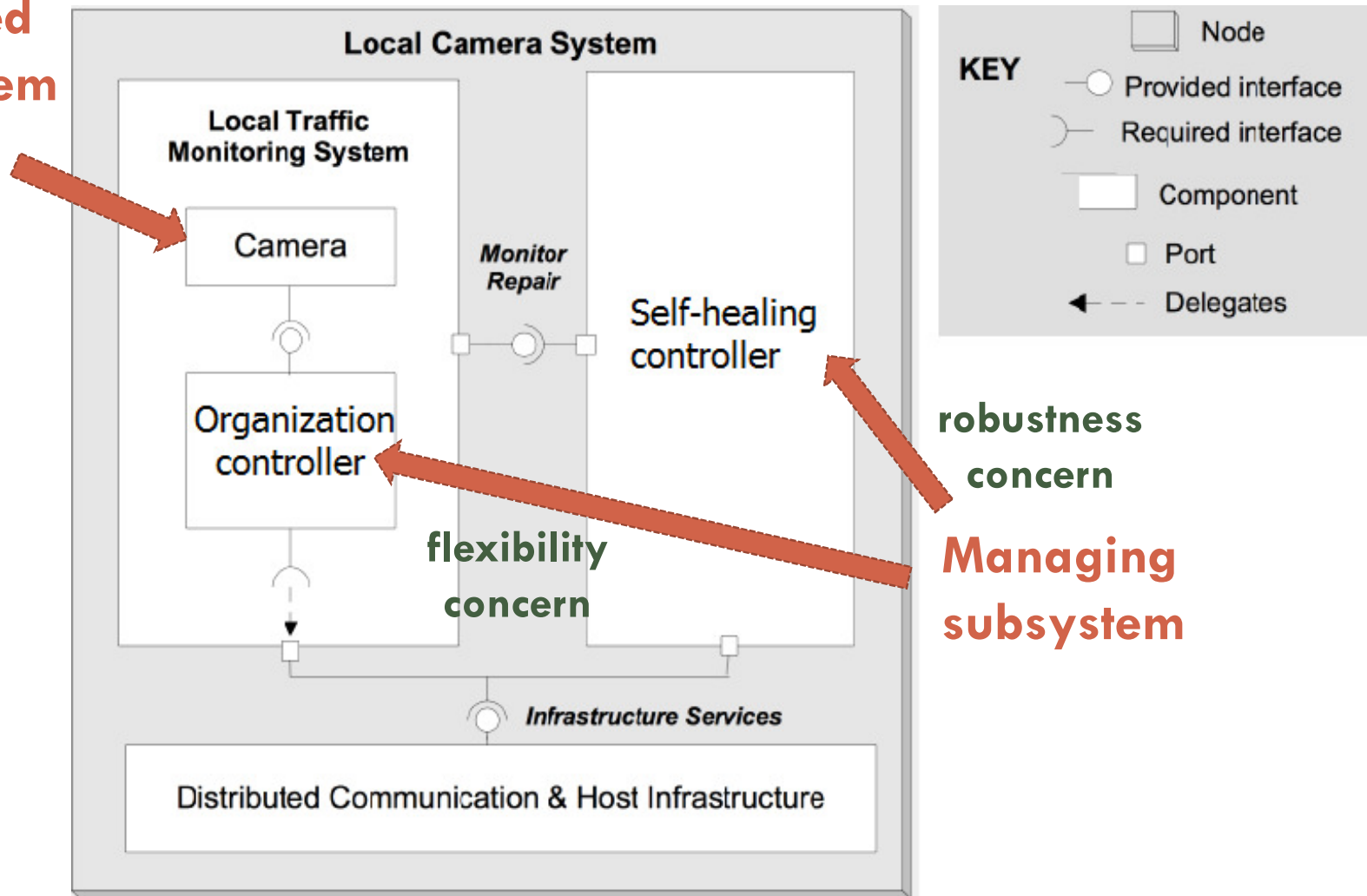
ROBUSTNESS

adaptation
concern

* M. U. Iftikhar and D. Weyns. A case study on formal verification of self-adaptive behaviors in a decentralized system. In FOCLASA 2012, Newcastle, U. K.

Running case study: camera system architecture

**Managed
subsystem**



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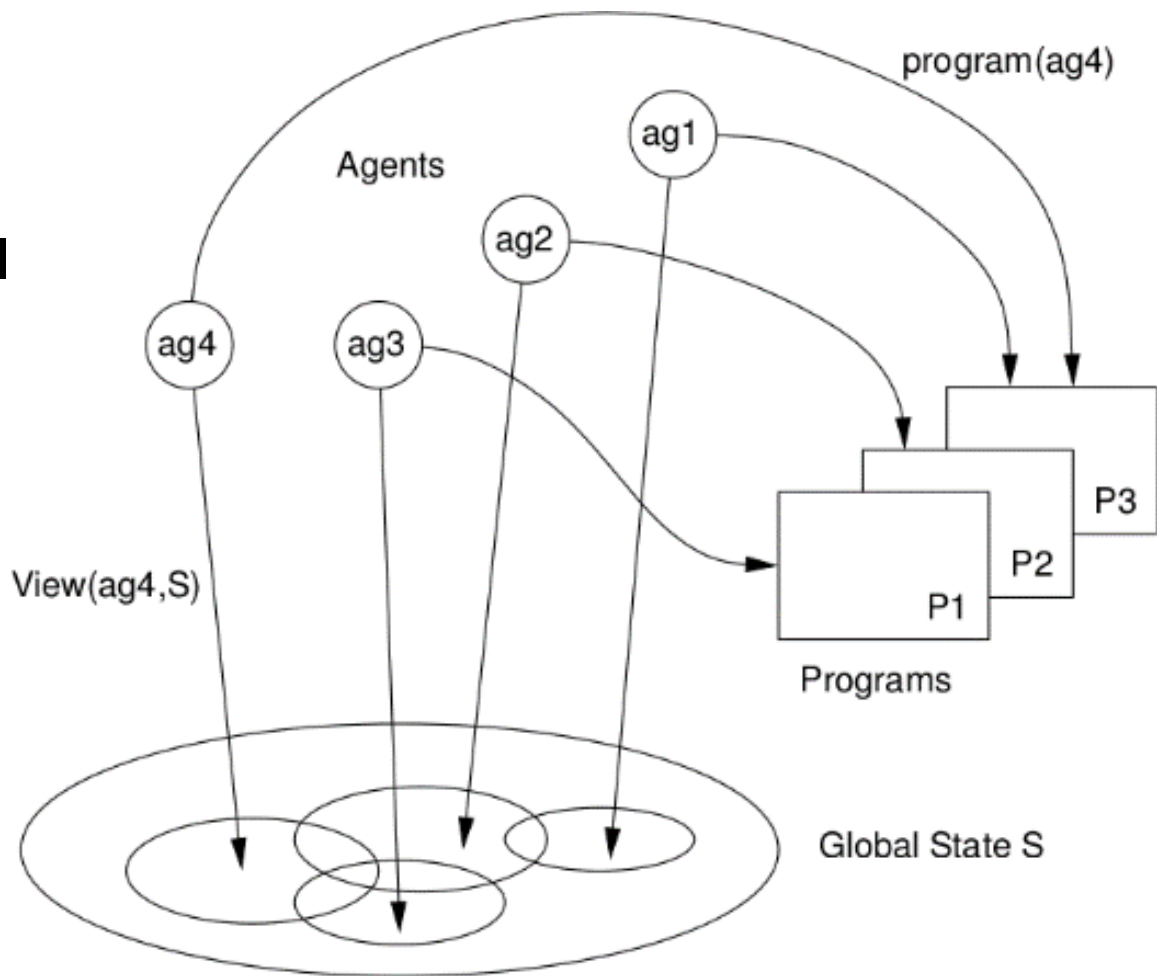
Abstract State Machines (ASMs)

- ▶ **ASMs are an extension of FSMs**
 - **states**: *multi-sorted first-order structures*, i.e. domains of objects with functions defined on them
 - **transitions**: named *transition rules* describing how functions change from one state to the next
- ▶ **Basic transition rule**: *if Condition then Updates*
where *Updates* is a set of **function updates** $f(t_1, \dots, t_n) := t$ **simultaneously executed when Condition** is true
- ▶ More complex **rule constructors** exist:
 - parallel (**par**) and sequential actions (**seq**)
 - non-determinism (**choose**)
 - unrestricted synch. parallelism (**forall**)
 - etc.

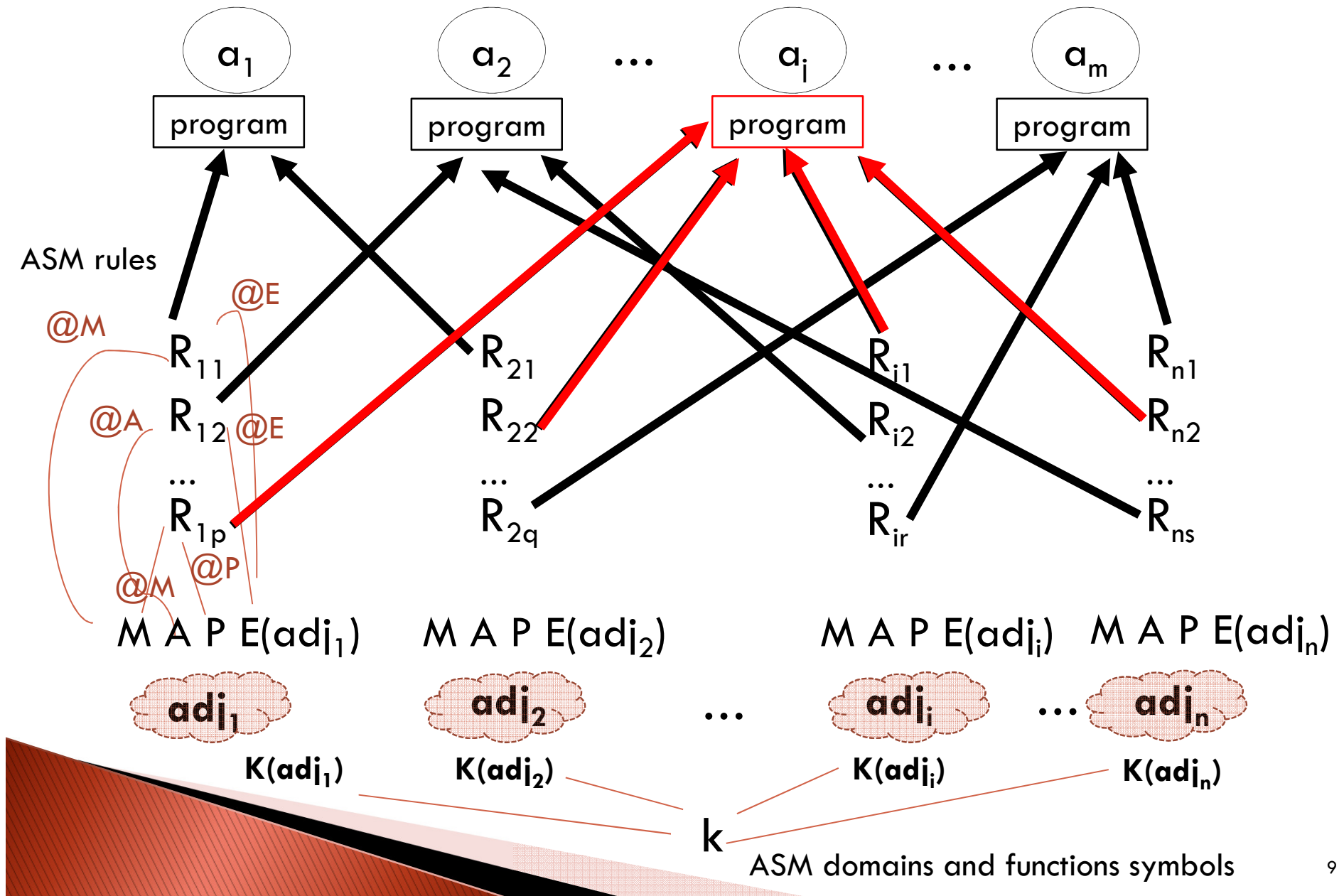
Multi-agent ASM (or distributed ASM)

Each **agent** $a \in \mathbf{AGENT}$

- ▶ has a “local” view $View(a, S)$ of the global state S
- ▶ executes its own program $prog(a)$ (i.e., an ASM rule) to determine the next global state



ASM model topology of the managing layer



Self-adaptive ASM

A multi-agent ASM:

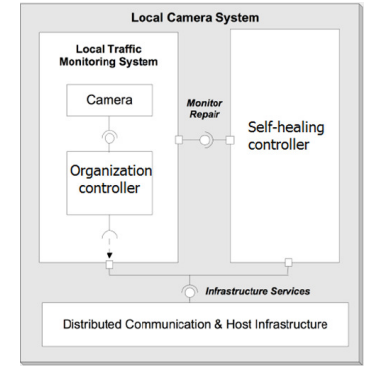
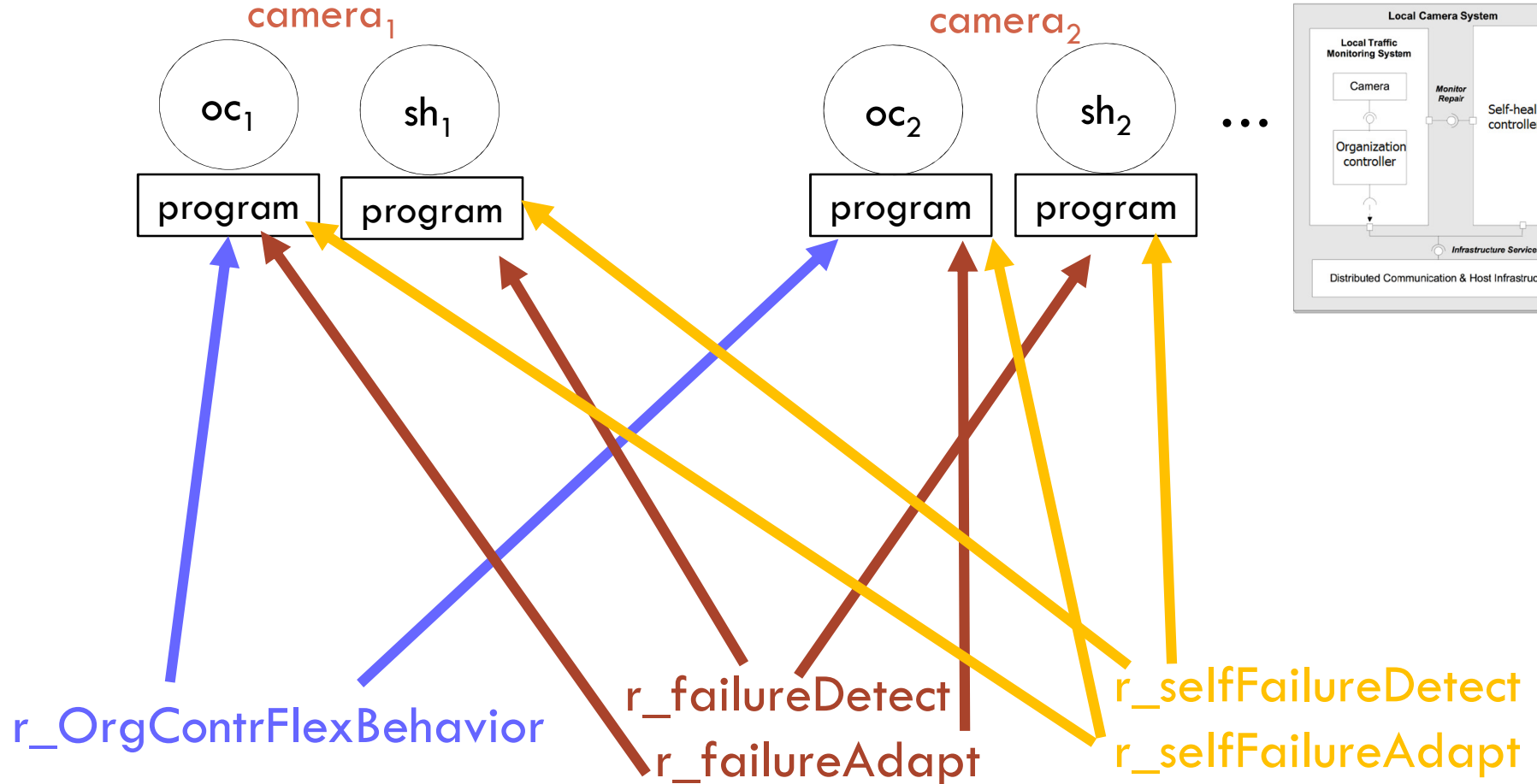
- ▶ **managed agents** $MdA \subseteq Agent$ encapsulate the system's **functional logic**
- ▶ **managing agents** $MgA \subseteq Agent$ encapsulate the **adaptation logic of MAPE-K loops**
- ▶ A common **knowledge** $K = \bigcup_{adj} K(adj)$ is shared by all managing agents
- ▶ The notion of **environment** is represented by **ASM monitored functions**
- ▶ A **MAPE loop** for an **adaptation concern** adj_i :

$$MAPE(adj_i) = \{R_{MAPE(adj_i)}^{a_1}, \dots, R_{MAPE(adj_i)}^{a_m}\}$$

- $\{a_1 \dots, a_m\} \subseteq MgA$ are the managing agents involved in the loop
 - $R_{MAPE(adj_i)}^{a_j}$ is the behavioral contribution of a_j to the loop
- ▶ The **program of a managing agent** a_j is the parallel execution of all its behavioral contributions to the loops j_1, \dots, j_k it is involved to:

$$program(a_j) = \mathbf{par} R_{MAPE(adj_{j_1})}^{a_j}, \dots, R_{MAPE(adj_{j_k})}^{a_j} \mathbf{endpar}$$

ASM model topology of the Traffic monitoring case study



MAPE(flexibility)
flexibility

MAPE (extFailure)
extFailure

MAPE(intFailure)
intFailure

k

Traffic monitoring case study

Program of each organization controller

```
macro rule r_organizationController =  
  par  
    orgContrFlexBehavior(self) //Adaptation due to congestion  
    r_failureAdapt[] //Adaptation due to external failure  
    r_selfFailureAdapt[] //Adaptation due to internal failure  
  endpar  
  
agent OrganizationController : r_organizationController[]
```

Excerpt of rule with MAPE computations

```
macro rule r_selfFailureAdapt =  
  par  
    if stopCam(camera(self)) then //@M_s  
      if state(camera(self)) != FAILED then //@A  
        state(camera(self)) := FAILED //@E  
      endif  
    endif  
    if startCam(camera(self)) then //@M_s  
      if state(camera(self)) = FAILED then //@A  
        par //@E  
          state(camera(self)) := MASTER  
          ...  
        endpar  
      endif  
    endif  
  endpar
```

Centralized self-aware monitoring

if *Cond* **then** *Analyze* //@M_c[s]

ASM rule schemes or patterns capture the general semantics of MAPE computations

Formal analysis techniques

Supported by the toolset **ASMETA** (*ASM mETAmodeling*)

▶ **Model validation**

- provide early feedback, less demanding than property verification
- Techniques
 - **Simulation** (interactive simulation, random simulation)
 - **Scenario-based validation**

▶ **Model verification**

- based on the *model checking* technique
 - **Model review:** verification of *meta-properties* (system-independent properties) defined as CTL formulae
 - **Verification of invariants and adaptation goals** expressed in CTL/LTL formulas

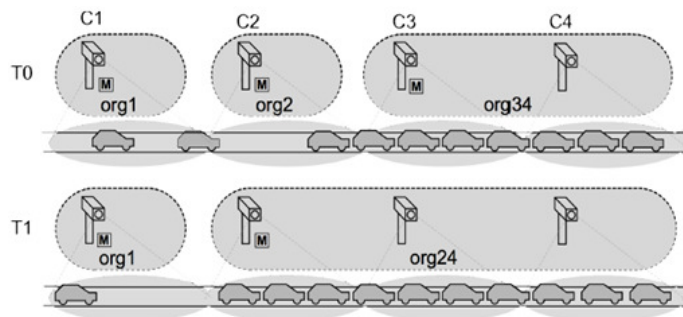
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Scenario-based validation

- ▶ Definition of key scenarios specifying the expected behavior of the model
- ▶ Scenarios are written in the language Avalla and
- ▶ executed through the validator ASMETA/AsmetaV

Example

Flexibility scenario from T0 to T1 in Avalla



```
scenario Flexibility_T0_T1
load main.asm
```

```
set stopCam(c1) := false; set stopCam(c2) := false; set stopCam(c3) := false;
set stopCam(c4) := false; set startCam(c1) := false; set startCam(c2) := false;
set startCam(c3) := false; set startCam(c4) := false; set congestion(c1) := false;
set congestion(c2) := false; set congestion(c3) := true; set congestion(c4) := true;
set elapsedWaitTime(shc3) := false; set elapsedWaitTimePlusDelta(shc4) := false;
exec par
```

```
state(c3) := MASTERWITHSLAVES
state(c4) := SLAVE
slaves(c3, c4) := true
getMaster(c4) := c3
congested(oc3) := true
congested(oc4) := true
```

```
endpar;
```

```
step
```

```
set congestion(c2) := true;
```

```
step
```

```
check getMaster(c4)=c3 and s_offer(c3)=true and s_offer(c4)=false and
slaves(c3,c4)=true and state(c1)=MASTER and state(c2) = MASTER and
state(c3) = MASTERWITHSLAVES and state(c4)=SLAVE;
```

```
step
```

```
check isAlive(c4)=false and newSlave(c2,c3)=true and getMaster(c4)=c3 and
s_offer(c3)=true and s_offer(c4)=false and slaves(c3,c4)=false and
state(c1)=MASTER and state(c2) = MASTER and state(c3) = SLAVE and
state(c4)=SLAVE;
```

```
step
```

```
check isAlive(c4) = false and newSlave(c2,c3) = false and getMaster(c4) = c3 and
s_offer(c3) = true and s_offer(c4) = false and slaves(c2,c3) = true and
slaves(c2,c4) = true and state(c1) = MASTER and
state(c2) = MASTERWITHSLAVES and state(c3) = SLAVE and
state(c4) = SLAVE;
```

Model verification

through the ASMETA/AsmetaSMV that translates ASM into models of the model checker NuSMV

▶ Invariant verification:

I1: $\text{ag}(\text{not}(\text{forall } \$c \text{ in Camera with state}(\$c) = \text{SLAVE}))$

I2: $\text{ag}(\text{not}(\text{forall } \$c \text{ in Camera with state}(\$c) = \text{MASTERWITHSLAVES}))$

▶ Adaptation goals:

Flexibility **F1:** $\text{ag}((\text{state}(c_i) = \text{MASTER and congested}(oc_i) \text{ and state}(c_{i+1}) = \text{MASTER and congested}(oc_{i+1})) \text{ implies af}(\text{state}(c_i) = \text{MASTERWITHSLAVES and slaves}(c_i, c_{i+1})))$

Robustness **R1:** $\text{ag}((\text{stateC}(c_i) = \text{FAILED and slaves}(c_i, c_{i+1})) \text{ implies ef}(\text{not}(\text{slaves}(c_i, c_{i+1}))))$

Model review

- ▶ through the AsmetaMA tool (based on AsmetaSMV)
- ▶ a meta-property violation may indicate the presence of a real fault or only of a stylistic defect
- ▶ **Meta-properties categories** for SA:
 - **MPnc: MAPE loops are not in conflict.** discover unwanted interferences between MAPE-K loops in terms of inconsistent ASM function updates
 - **MPe : all rules involved in MAPE loops are executed,** i.e., there is no over specification inside a MAPE loop
 - **MPm: the knowledge is minimal,** i.e., it does not contain locations that are unnecessary

Faced challenges

- ▶ **Formal modeling self-adaptive behavior through a clear separation of concerns in a decentralized view**
 - By distinguishing ASM managing agents from managed ones
 - By identifying different adaptation concerns
 - By distributing the MAPE computations of a loop among agents
 - By treating, inside the behavior of a managing agent, different adaptation concerns
 - By distinguishing between decentralized and centralized loop's control through specific ASM rule patterns
- ▶ **Formal functional analysis**
 - Validate adaptation requirements by simulation
 - Determine conflicting MAPE loops
 - Assert the system correctness by model checking a set of properties expressing invariants and adaptation goals
 - Check for model completeness without overspecification

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Conclusion and future work

- ▶ **Self-adaptive ASMs** allowed us to model and analyze the behavior of self-adaptive systems formally
 - in terms of MAPE-K control loops executed by ASM agents
- ▶ **Validation and verification** techniques allowed us to ensure the **functional correctness of the adaptation logic** by discovering interfering adaptation concerns and goals
- ▶ In the future, we want to exploit runtime monitoring techniques for runtime verification
- ▶ We also want to exploit extensions of ASMs with time models for specifying time-triggered adaptation



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